Adding Persistence to Java

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Overview

- One possible approach to adding persistence to the Java platform:
  - Rationale
  - Model
  - Implementation
  - Open questions
• Widely used

Goal: minimize disruption of current users’ knowledge and practices

• Lots of existing code

Goal: maximize ability to reuse existing code
How is persistent state expressed and accessed?
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Could use:

- Special constructor for persistent objects (`pnew`)
- Inherit from a special `PersistentObject` class
- Add `@Persistent` annotation on a class
How is persistent state expressed and accessed?

Could use:

- Special constructor for persistent objects ($pnew$)
- Inherit from a special PersistentObject class
- Add @Persistent annotation on a class

All rejected as being non-orthogonal — prevents the use of a volatile class/object for persistence and vice versa.
What is the model for consistency and recovery?
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Could add transactions at the object level; nested transactions with compensation; transaction-based concurrency control.
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Could add transactions at the object level; nested transactions with compensation; transaction-based concurrency control

Rejected as too invasive and disruptive to existing code and practices
Persistence by reachability

- Designate some variables as persistence roots
- All objects reachable from a persistent root are persistent
- Lots of research on this from ~1980—2000, including an implementation for Java at Sun Labs (known initially as OPJ, later as PJama).

However...
Hard state vs soft state

**Hard state** is information that cannot be lost

**Soft state** is information that can be reconstructed from existing hard state

From *Manageability, availability and performance in Porcupine: a highly scalable, cluster-based mail service*, Saito, Bershad, and Levy, 1999
What is this man about to say?
What is this man about to say?

“Have you tried turning it off and on again?”

https://www.youtube.com/watch?v=nn2FB1P_Mn8
Hard & soft state in NVRAM

Soft state:

- *Could* discard at uncontrolled restarts
- Or, discard *only* when corruption is suspected

Hard state:

- *Cannot* discard. *Must* provide backup, recovery, etc., perhaps replication (for availability)
Our model

• Partitioned heap (following NVM Direct, see Bill Bridge’s presentation tomorrow)

  • Allows enforced separation of state

• Partition checkpoints
Partitioned heaps

- Partition the heap into volatile and persistent *regions*, visible to the programmer.

- Persistent regions are mapped from DAX files.

- Each object is allocated in a specific region.

- Disallow cross-region references, except from volatile regions.

- When `main` begins, there is a single volatile region.
volatile

main

P2

P1

Persistent

13
The current region

- Each thread has a notion of *the current region* — all allocations by the thread are in the current region.

- The current region can be changed. Code can be wrapped to use a specific region, e.g.:

  ```java
  myRegion.run(() -> {h = new HashMap();});
  ```

- The persistence of a data structure can be specified by context—enables reuse.
Checkpoints

- The program can invoke a checkpoint primitive which snapshots the current state of one or more regions.

- At re-attach, a region is recovered to the last checkpoint.

- Regions can be checkpointed independently or together. Example: checkpoint a DB infrequently, but the log after every addition.
Epochs

• Each region goes through a series of epochs. An epoch is the period between two checkpoints (or before the first).

• A checkpoint freezes the state of the region within that epoch. The data within an old epoch are immutable.

• When an object is modified for the first time within the current epoch, a copy is made. All subsequent modifications within the current epoch are to this copy.

• The running program can only observe the most recent version of an object.
Recovery

• When a restart occurs, and the latest epoch is uncommitted, recovery takes place.

• Recovery discards the current, uncommitted epoch and reverts to the state at the end of the previous epoch.

• Locks held on objects in the region are re-initialized.
Managing checkpoints

- Each checkpoint persists until explicitly discarded.

- If the oldest checkpoint is deleted, objects in that checkpoint without modified copies in the next epoch are absorbed into the next epoch. If an object is superseded by a copy in the next epoch, it can be discarded.

- An intermediate-age checkpoint can also be deleted, merging the two epochs around it.

- We can also recover to any checkpoint, discarding all the later epochs.
The programmer’s burden

To convert an existing application to use NVM, the programmer:

- Partitions the heap into volatile and non-volatile regions
- Adds the region creation and loading logic,
- Wraps non-volatile object creations in Region.run().
- Adds calls to checkpoint() at places where the region is consistent.
Example: simple phone directory
(following an example by Eliot Moss)

class PhoneDirectory {
    HashMap<String, String> dir;

    public static void main(String[] args) {
        ...
    }
    ...
}
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```java
public class PhoneDirectory extends Region {
    HashMap<String, String> dir;

    public static void main(String[] args) {
        ...
    }
    ...
}
```
public class PhoneDirectory extends Region {

    HashMap<String, String> dir;

    public static void main(String[] args) {
        PhoneDirectory pd = null;
        try {
            pd = new PhoneDirectory("phonedir.region");
        } catch (InvalidRegionFileException | IOException ex) {
            System.err.println("Cannot connect/create region");
            System.exit(1);
        }
    }
}

...
public class PhoneDirectory extends Region {

    HashMap<String, String> dir;

    public static void main(String[] args) {
        PhoneDirectory pd = null;
        try {
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        } catch (InvalidRegionFileException | IOException ex) {
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            System.exit(1);
        }
    ...

    }

    @Override
    public void initialize() { this.dir = new HashMap<>(); }
}
Further along in main():

// command received to add an entry for name and number
...
addEntry(name, number);
Further along in main():

// command received to add an entry for name and number
...
run(() -> addEntry(name, number));

Region management added
private void addEntry(String name, String number) throws PDEException {
    if (dir.containsKey(name)) {
        error("name already in directory");
    }
    dir.put(name, number);
}

Without persistence
private void addEntry(String name, String number) throws PDException
{
    if (dir.containsKey(name)) {
        error("name already in directory");
    }
    dir.put(name, number);
    checkpoint();
}
private void addEntry(String name, String number) throws PDException {
    String name2 = copyString(name);
    if (dir.containsKey(name2)) {
        error("name already in directory");
    }
    dir.put(name2, copyString(number));
    checkpoint();
}
private void addEntry(String name, String number) throws PDException {
    String name2 = copyString(name);
    if (dir.containsKey(name2)) {
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- The String parameters must be copied to the region.
private void addEntry(String name, String number) throws PDException {
    String name2 = copyString(name);
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        error("name already in directory");
    }
    dir.put(name2, copyString(number));
    checkpoint();
}

• The String parameters must be copied to the region.

• For immutable data, the JVM could do this automatically.
@CopyArgs
private void addEntry(String name, String number) throws PDException {
    if (dir.containsKey(name)) {
        error("name already in directory");
    }
    dir.put(name, number);
    checkpoint();
}

• Using annotations to save boilerplate
Implementation

• A region can be mapped anywhere (although it may be possible to have it commonly mapped at the same address). Hence it must be position-independent, or cheaply (and preferably incrementally) relocatable.

• We’d like checkpoints to be fast, so they can be used relatively often

• We’d like checkpoints of small regions to be even faster

• Steady-state performance of long-running programs is paramount

• Crashes and recovery are infrequent

• NB: NVM is much cheaper (per byte) than DRAM, so we can trade space for time—but still be mindful of cache behaviors
The heap structure of a region

- An object ID is a reference to an Object Table Entry (OTE)
  - Self-relative when in the heap
- The OT references the most recent version of an object.
- Each version references the next oldest from its header (reserved header space).
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Open questions

1. Is this a useful model? When is it not usable? What kinds of mistakes are made, and what is the mitigation?

2. Which libraries and apps need to be changed? How do we find them?

3. How to handle threads? Objects with external soft state?

4. Make it work; make it fast. How fast?
   - Peak performance, checkpoint latency

5. How to enable software evolution?
Questions?
Comments?
Integrated Cloud
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